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Reg No.:

Name:

APJ ABDUL KALAM TECHNOLOGICAL UNIVERSI

Fifth semester B.Tech degree examinations (S) September 2020

Course Code: CS301 Course Name: THEORY OF COMPUTATION

Max. Marks: 100

PART A

Duration: 3 Hours

ages

Answer all questions, each carries3 marks.

Marks

Formally define extended delta for an NFA. Show the processing of input (3) w = 0101 for the following NFA.



- 2 Differentiate between the transition function in DFA, NFA and ϵ -NFA (3)
- 3 Design a Moore machine to determine the residue of mod 2 of the input treated (3) as a binary string.
- 4 Give a regular expression for the set of all strings not containing 101 as a (3) substring

PART B

Answer any two full questions, each carries 9 marks.

5 a) Convert the following NFA to DFA and describe the language it accepts. (5) $M = (\{P, Q, R, S, T\}, \{0,1\}, \delta, P, \{S, T\})$ and δ is given as:

	0	1
Р	{P,Q}	{P}
Q	{ R , S }	{T}
R	{P,R}	{ T }
S	-	-
Τ	-	

b) Prove that "A language L is accepted by some ϵ -NFA if and only if L is (4) accepted by some NFA"

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6 a) State Myhill-Nerode theorem, Minimize the following DFA.



- b) Find an equivalent ϵ -NFA for the following regular expression $(0 + 1)^* 011$
- 7 a) Convert the following ϵ -NFA to NFA

	E	1	2	3
q 0	Ø	{ q0}	{ q1}	{ q2}
q1	{ q0}	{ q1}	{ q2}	Ø
q2	{ q1}	{ q2}	Ø	{ q0}

b) Describe clearly the equivalent classes of the Canonical Myhill-Nerode relation (5) for the language of binary strings with second-last symbol as 0.

PART C Answer all questions, each carries3 marks.

8 State the closure properties of regular sets. (3)
9 Define context free grammar. Consider the following CFG (3)
S → aS | Sb |a | b
Prove by induction on the string length that no string in L(G) has ba as substring.
10 Design a PDA to accept the set of strings with twice as many 0's as 1's. (3)
11 List the decision problems related with type 3 Formalism. (3)

PART D

Answer any two full questions, each carries9 marks.

- 12 a) State pumping lemma for regular languages. Prove that the language L = (5) $\{a^{n^2} | n > 0\}$ is not regular.
 - b) Convert the following grammar into Chomsky normal form

 $S \rightarrow ASB | \epsilon, \quad A \rightarrow aAS | a, \quad B \rightarrow SbS | A | bb$

(5)

(4)

(4)

(4)

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13 a) Prove the equivalence of acceptance of a PDA by final state and empty of the		
b) Define a deterministic PDA. How a DPDA differs from a new local differs.		
PDA?	c (3)	
14 a) Let G be the grammar		
$S \rightarrow aB bA, A \rightarrow a aS bAA, B \rightarrow b bS aBB$	(4)	
For the string aabbaabbba find		
i) leftmost derivation, ii) parse tree, and iii) is the grammar ambiguage		
b) Design a PDA to accept the language $L = \{ww^R \mid w \in \{0, 1\}^*\}$		
PART E	(5)	
Answer any four full questions, each carries 10 marks		
15 a) Show that the language $L = \{ww w \in \{a, b\}^*\}$ is not a CFL.	(5)	
b) Design a TM to compute the 2's complement of a binary string.	(5)	
16 a) State and prove pumping lemma for context free languages. Mention the	(6)	
application of pumping lemma.	(0)	
b) Design a Turing machine to accept,	(4)	
$L = \{ w \in \{0,1\}^* w \text{ has equal number of } 0's \text{ and } 1's \}.$	()	
a) Compare context sensitive grammar and context free grammar. Can we design a	(5)	
PDA for context sensitive languages? Justify your answer.	(5)	
b) Design a TM to find the sum of two numbers m and n. Assume that initially the	(5)	
tape contains m number of 0s followed by # followed by n number of 0s 18° c) 4°	(0)	
Are there any languages which are not recursively enumerable, but accepted by	(5)	
a multi-tape Turing machine? Justify your answer.		
b) Define formally Type 0, Type 1, Type 2 and Type 3 grammar. Show the	(5)	
19 a) List the element of the second		
 b) Define a U in the closure properties of Recursive Languages 	(4)	
b) Define a Universal Turing Machine (UTM). With the help of suitable arguments	(6)	
20 a) Compare more in the simulation of other Turing machines by a UTM.		
b) Show that the above of	(3)	
c) Show that the class of recursive languages is closed under complementation.	(3)	
complementation	(4)	
comprementation.		
