

**APJ ABDULKALAM TECHNOLOGICAL UNIVERSITY  
08 PALAKKAD CLUSTER**

Q. P. Code : 2A181

(Pages: 4)

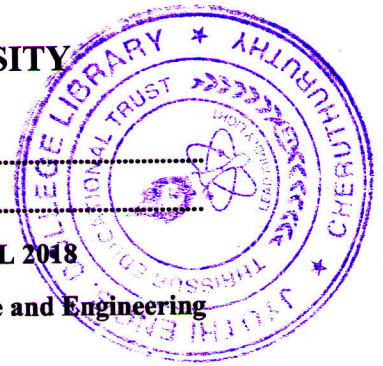
Name .....

Reg. No: .....

**SECOND SEMESTER M.TECH. DEGREE EXAMINATION, APRIL 2018**

Branch: CSE

**Specialization: Computer Science and Engineering**



**08 CS 6012 ADVANCED COMPILER DESIGN**

Time:3 hours

Max. Marks: 60

Answer all six questions.

Modules 1 to 6: Part 'a' of each question is compulsory and answer either part 'b' or part 'c' of each question.

Q.no.	Module 1	Marks												
1.a	Obtain an annotated parse tree for the following syntax directed definition with synthesised and inherited attributes for the grammar for type declarations.	3												
	<table border="0"> <tr> <td>Production</td> <td>Semantic rule</td> </tr> <tr> <td>D → T L</td> <td>L.in := T.type</td> </tr> <tr> <td>T → int</td> <td>T.type := integer</td> </tr> <tr> <td>T → real</td> <td>T.type := real</td> </tr> <tr> <td>L → L<sub>1</sub>, id</td> <td>L<sub>1</sub>.id := L.in; addtype (id.entry, L.in)</td> </tr> <tr> <td>L → id</td> <td>addtype(id.entry, L.in)</td> </tr> </table>	Production	Semantic rule	D → T L	L.in := T.type	T → int	T.type := integer	T → real	T.type := real	L → L <sub>1</sub> , id	L <sub>1</sub> .id := L.in; addtype (id.entry, L.in)	L → id	addtype(id.entry, L.in)	
Production	Semantic rule													
D → T L	L.in := T.type													
T → int	T.type := integer													
T → real	T.type := real													
L → L <sub>1</sub> , id	L <sub>1</sub> .id := L.in; addtype (id.entry, L.in)													
L → id	addtype(id.entry, L.in)													

Answer b or c

- |   |   |   |
|---|---|---|
| b | Construct a Syntax-Directed Translation scheme that translates arithmetic expressions from infix into postfix notation. Your solution should include the context-free grammar, the semantic attributes for each of the grammar symbols, and semantic rules. Show the application of your scheme to the input "3*4+5*2". | 6 |
| c | Explain parser stack implementation of Postfix SDT using an example of desk calculator.   | 6 |

**Module 2**

2.a	Consider following definition	3														
	<table border="0"> <tr> <td>D → T L</td> <td>L.in := T.type</td> </tr> <tr> <td>T → real</td> <td>T.type := real</td> </tr> <tr> <td>T → int</td> <td>T.type := int</td> </tr> <tr> <td>L → L<sub>1</sub>, I</td> <td>L<sub>1</sub>.in := L.in; I.in = L.in</td> </tr> <tr> <td>L → I</td> <td>I.in = L.in</td> </tr> <tr> <td>I → I<sub>1</sub> [num]</td> <td>I<sub>1</sub>.in = array(numeral, I.in)</td> </tr> <tr> <td>I → id</td> <td>addtype(id.entry, I.in)</td> </tr> </table>	D → T L	L.in := T.type	T → real	T.type := real	T → int	T.type := int	L → L <sub>1</sub> , I	L <sub>1</sub> .in := L.in; I.in = L.in	L → I	I.in = L.in	I → I <sub>1</sub> [num]	I <sub>1</sub> .in = array(numeral, I.in)	I → id	addtype(id.entry, I.in)	
D → T L	L.in := T.type															
T → real	T.type := real															
T → int	T.type := int															
L → L <sub>1</sub> , I	L <sub>1</sub> .in := L.in; I.in = L.in															
L → I	I.in = L.in															
I → I <sub>1</sub> [num]	I <sub>1</sub> .in = array(numeral, I.in)															
I → id	addtype(id.entry, I.in)															

Obtain the parse tree and dependency graph for the string `int x[3], y[5]`

**Answer b or c**

**b** Obtain three address codes for the following statement and explain. 6

```

while a < b do
  if c < d then
    x := y + z
  else
    x := y - z

```

**c** Consider grammar and rules given below for array address translation and generating address code for array references: 6

```

E → E1 + E2           {E.addr = newtemp();
                        gen(E.addr '=' E1.addr '+' E2.addr);}

E → E1 * E2           {E.addr = newtemp();
                        gen(E.addr '=' E1.addr '*' E2.addr);}

E → id                {E.addr = id.lexeme; }

E → L                 { E.addr = newtemp();
                        gen(E.addr '=' L.array.basename '['L.addr ']'); }

L → id [ E ]          {L.array = id.lexeme;
                        L.type = L.array.typeofelement;
                        L.addr=newtemp();
                        gen(L.addr '=' E.addr '*' L.type.width);}

L → L1 [ E ]          {L.array = L1.array;
                        L.type=L1.type.typeofelement;
                        t = newtemp(); L.addr = newtemp();
                        gen(t '=' E.addr '*' L.type.width);
                        gen(L.addr '=' L1.addr '+' t); }

```

Function newtemp() returns a new temporary name  
L.array.basename means name of the array  
L.array.typeofelement means type of the element of the array  
L.type.width means width of L.type

Assume size of integer to be 4 bytes, and lower bound of the arrays to be 0

Let A, B and C be 3X4, 4X5, and 3X5 arrays of integers respectively. Let i, j, and k be integers.

Construct an annotated parse tree for the expression C[i][j] + A[i][k]\*B[k][j] and show the 3-address code sequence generated for the expression.

Q.no.	Module 3	Marks
3.a	What is activation record?	3

**Answer b or c**

**b** Write the algorithm for Mark-and-Compact Garbage Collector and explain. 6

- c Explain the use of access links for finding non local data. Also explain the display implementation. 6

**Q.no.** **Module 4** **Marks**  
4.a Write the techniques for the optimization of basic blocks. 3

**Answer b or c**

- b Apply simple code generation algorithm to generate code for the following statements. 6

```
t1=a-b
t2=c+d
c=e*t2
d=t1+t2
```

Show the instructions generated and changes in address and register descriptors while generating code each statement. Assume there are only three registers .t1 and t2 are temporary variables.

- c Generate code for the following statements from labelled expression tree using Ershov numbers. t,u,v, and w are temporary variables . Assume the number of registers available is three. 6

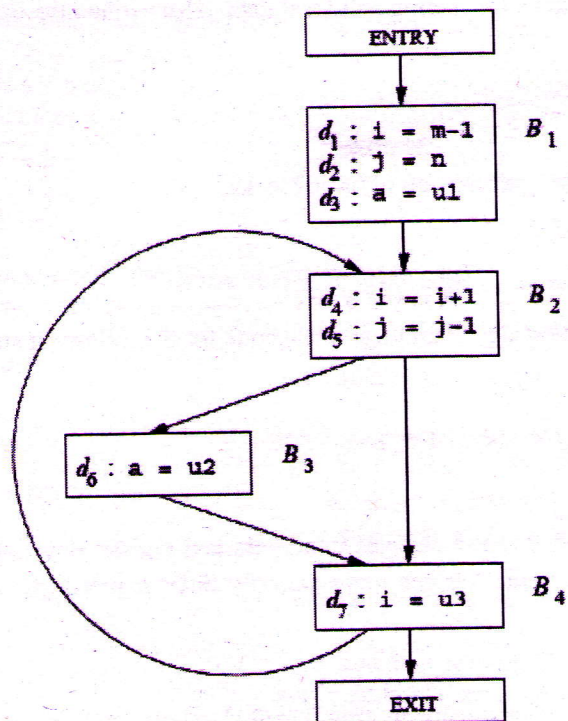
```
t= a-b
u=c+d
v=e*u
w=t+v
```

**Q.no.** **Module 5** **Marks**  
5.a What are dominators? Draw a flow graph with at least one loop and obtain the dominators of each node. 4

**Answer b or c**

- b For the flow graph shown below perform live variable analysis and explain 8





c Write the algorithm for finding the available expression and explain with the help of a suitable example 8

<b>Q.no.</b>	<b>Module 6</b>	<b>Marks</b>
--------------	-----------------	--------------

6.a	Explain true dependence and name dependence with the help of suitable examples.	4
-----	---	---

**Answer b or c**

b	Explain global code motion in detail	8
---	--------------------------------------	---

c	Explain the concept of loop unrolling with the help of suitable examples.	8
---	---	---